# Lecture 14: DNA Sequencing 

Study Chapter 8.9
Midterm on Tuesday 10/15
Open book, open notes, no computer
Study Session on 10/14 in FB008 from 5pm-7pm

## DNA Sequencing



- Shear DNA into millions of small fragments
- Read 500-700 nucleotides at a time from the small
fragments
(Sanger method)



## Fragment Assembly



- Assembles the individual overlapping short fragments (reads) into a genomic sequence
- Shortest Superstring problem from last time is an overly simplified abstraction
- Problems:
- DNA read error rate of 1\% to 3\%
- Can't separate strands
- DNA is full of repeats
- Let's take a closer look



## Traditional DNA Sequencing

 DNA


## Different Types of Vectors



| VECTOR | $\underline{\text { Size of insert (bp) }}$ |
| :---: | :---: |
| Plasmid | $2,000-10,000$ |
| Cosmid | 40,000 |
| BAC (Bacterial Artificial <br> Chromosome) | $70,000-300,000$ |
| YAC (Yeast Artificial <br> Chromosome) | $>300,000$ <br> Not used much <br> recently |



## Dideoxy (Sanger) Sequencing



## Template strand - g t a a g a c t g t

Coding strand - c a t t c t g a c a
ddT Reaction -

ddC Reaction -

ddG Reaction -

-ddA Reaction -


## Challenging to Read Answers

 Electropherogram


## Reading an Electropherogram



- Issues
- Noisy start up due to anomalous migration of short fragments that carry bulky dyes
- Traces become less uniform as run proceeds
- Large dye responses can overwhelm succeeding lower amplitude responses
- Occasional mismatches of reaction with template
- Methods for calling the nucleotides: PHRED
- Base calls
- Maintains quality scores
- Monitors peak positions


## Shotgun Sequencing

 genomic segment


## Fragment Assembly




Cover region with $\sim 7$-fold redundancy
Overlap reads and extend to reconstruct the original genomic region

## Read Coverage




Length of genomic segment: L
Number of reads:
$n \quad$ Coverage $C=n l / L$
Length of each read:
$l$
How much coverage is enough?
Lander-Waterman model:
Assuming uniform distribution of reads, $C=10$ results in 1 gapped region per 1,000,000 nucleotides

## Challenges in Fragment Assembly



- Repeats: A major problem for fragment assembly
- $>50 \%$ of human genome is repeats:
- over 1 million Alu repeats (about 300 bp )
- about 200,000 LINE repeats (1000 bp and longer)



## Types of Genome Assemblies



- De Novo -

An assembly based entirely on self-consitency or self-similarity of short reads (contigs).

- Comparative Refers an assembly of a genome using the sequence of a close relative as a scaffold or reference. Sometimes called a "template assembly" or "a resequencing"
- Confounding problem for both types: Repeats


## Repeat Types



- Low-Complexity DNA (e.g. ATATATATACATA...)
- Microsatellite repeats $\quad\left(\mathrm{a}_{1} \ldots \mathrm{a}_{\mathrm{k}}\right)^{\mathrm{N}}$ where $\mathrm{k} \sim 3-6$ (e.g. CAGCAGTAGCAGCACCAG)
- Transposons/retrotransposons
- SINE
- LINE

Short Interspersed Nuclear Elements (e.g., Alu: ~300 bp long, >10 ${ }^{6}$ in human)

Long Interspersed Nuclear Elements $\sim 500-5,000$ bp long, > 200,000 in human

- LTR retroposons
- Gene Families
- Segmental duplications


## Overlap-Layout-Consensus


Assembler programs ARACHNE, PHRAP, CAP, TIGR, CELERA
Common Approach:
Overlap: find potentially overlapping reads

Layout: merge reads into contigs and then combine contigs into supercontigs


Consensus: requires many overlapping reads to derive the DNA seq..ACGATTACAATAGGTT.. uence and to correct for read errors

## Overlap



- Find the best match between the suffix of one read and the prefix of another (shortest superstring)
- Due to sequencing errors, most algorithms use dynamic programming to find the optimal overlap alignment
- Filter out fragment pairs that do not share a significantly long common substring


## Overlapping Reads



- Make an index of all $k$-mers of all reads

$$
(k \sim 20-24)
$$

- Find read-pairs sharing a k-mer
- Extend alignment -
throw away if not $>95 \%$ similar
tacA tagattacacagattact ga

TAGT TAGATTACACAGATTACTAGA


## Histogram Example



$$
v=\text { tagattacacagattattga }
$$

- Histogram of 3-mers (18 total)

|  | $\mathrm{A}_{2}$ | $\mathrm{C}_{2}$ | $\mathrm{G}_{2}$ | $\mathrm{~T}_{2}$ |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathrm{~A}_{3}: \mathrm{C}_{3}: \mathrm{G}_{3}: \mathrm{T}_{3}$ | $\mathrm{~A}_{3}: \mathrm{C}_{3}: \mathrm{G}_{3}: \mathrm{T}_{3}$ | $\mathrm{~A}_{3}: \mathrm{C}_{3}: \mathrm{G}_{3}: \mathrm{T}_{3}$ | $\mathrm{~A}_{3}: \mathrm{C}_{3}: \mathrm{G}_{3}: \mathrm{T}_{3}$ |
| $\mathrm{~A}_{1}$ | $0: 0: 0: 0$ | $2: 0: 0: 0$ | $2: 0: 0: 0$ | $0: 0: 0: 3$ |
| $\mathrm{C}_{1}$ | $0: 1: 1: 0$ | $0: 0: 0: 0$ | $0: 0: 0: 0$ | $0: 0: 0: 0$ |
| $\mathrm{G}_{1}$ | $0: 0: 0: 2$ | $0: 0: 0: 0$ | $0: 0: 0: 0$ | $0: 0: 0: 0$ |
| $\mathrm{~T}_{1}$ | $0: 1: 1: 1$ | $0: 0: 0: 0$ | $1: 0: 0: 0$ | $2: 0: 1: 0$ |

## Overlapping Reads and Repeats



- Does this really speed up the process?
- A $k$-mer that appears N times, initiates $\mathrm{N}^{2}$ comparisons (you consider all pairs ofs reads that share the $k$-mer substring)
- For an Alu that appears $10^{6}$ times $\rightarrow 10^{12}$ comparisons - too much
- How to avoid repeats:

Discard all $k$-mers that appear more than

$$
t \times \text { Coverage, }(t \sim 10)
$$

## Finding Overlapping Reads



# k-mer table makes it easy to create local multiple alignments from the overlapping reads 



## Finding Overlapping Reads (cont'd)



- Correct errors using multiple alignment

- Score alignments
- Accept alignments with good scores


## Layout



- Repeats are still a major challenge
- Do two aligned fragments really overlap, or are they from two copies of a repeat?
- Solution: repeat masking - hide the repeats!!!
- Masking results in high rate of misassembly (up to 20\%)
- Misassembly means alot more work at the finishing step


## 2. Merge Reads into Contigs



- Overlap graph:
- Nodes: reads $r_{1} \ldots . r_{n}$
- Edges: overlaps ( $\mathrm{r}_{\mathrm{i}}, \mathrm{r}_{\mathrm{j}}$, shift, orientation, score)


Reads that come from two regions of the genome (blue and red) that contain the same repeat


## 2. Merge Reads into Contigs




We want to merge reads up to potential repeat boundaries

## 2. Merge Reads into Contigs




- Ignore non-maximal reads
- Merge only maximal reads into contigs


## 2. Merge Reads into Contigs




- Remove transitively inferable overlaps
- If read r overlaps to the right reads $r_{1}, r_{2}$, and $r_{1}$ overlaps $r_{2}$, then ( $r, r_{2}$ ) can be inferred by ( $\mathrm{r}, \mathrm{r}_{1}$ ) and ( $\mathrm{r}_{1}, \mathrm{r}_{2}$ )



## 2. Merge Reads into Contigs




## 2. Merge Reads into Contigs




- Ignore "hanging" reads, when detecting repeat boundaries


## Overlap graph after forming <br> Dalinelion Contios

 Target

Fragments


Fall 2013

Unitigs:
Gene Myers, 95

## Repeats, errors, and contig lengths

$D$ IIITM自

- Repeats shorter than read length are easily resolved
- Read that spans across a repeat disambiguates order of flanking regions
- Repeats with more base pair diffs than sequencing error rate are OK
- We throw overlaps between two reads in different copies of the repeat
- To make the genome appear less repetitive, try to:
- Increase read length
- Decrease sequencing error rate


## Role of error correction:

Discards up to $98 \%$ of single-letter sequencing errors
decreases error rate
$\Rightarrow$ decreases effective repeat content
$\Rightarrow$ increases contig length

## 2. Merge Reads into Contigs




- Insert non-maximal reads whenever unambiguous


## Link Contigs into Supercontigs




Normal density

Too dense:
Overcollapsed?

Inconsistent links:
Overcollapsed?

## Link Contigs into Supercontigs $\left(c^{2} t^{\prime} d\right)$ <br> 

Find all links between unique contigs
Connect contigs incrementally, if $\geq 2$ links


## Link Contigs into Supercontigs (cont'd) <br> 

Fill gaps in supercontigs with paths of overcollapsed contigs


## Link Contigs into Supercontigs

(cont'd)
 d (A, B )
Contig A

Contig B
Define $G=(V, E)$

> V := contigs
$E:=(A, B)$ such that $d(A, B)<C$

Reason to do so: Efficiency; full shortest paths cannot be computed

## Link Contigs into Supercontigs

 (cont'd)


Define $T$ : contigs linked to either $A$ or $B$
Fill gap between $A$ and $B$ if there is a path in
G passing only from contigs in $T$

## Consensus



- A consensus sequence is derived from a profile of the assembled fragments
- A sufficient number of reads is required to ensure a statistically significant consensus
- Reading errors are corrected


## Derive Consensus Sequence



TAGATTACACAGATTACTGA TTGATGGCGTAA CTA TAGATTACACAGATTACTGACTTGATGGCGTAAACTA TAG TTACACAGATTATTGACTTCATGGCGTAA CTA TAGATTACACAGATTACTGACTTGATGGCGTAA CTA TAGATTACACAGATTACTGACTTGATGGGGTAA CTA


TAGATTACACAGATTACTGACTTGATGGCGTAA CTA

Derive multiple alignment from pairwise read alignments

## Derive each consensus base by weighted voting

## Some Assemblers



- PHRAP
- Early assembler, widely used, good model of read errors
- Overlap $\mathrm{O}\left(\mathrm{n}^{2}\right) \rightarrow$ layout (no mate pairs) $\rightarrow$ consensus
- Celera
- First assembler to handle large genomes (fly, human, mouse)
- Overlap $\rightarrow$ layout $\rightarrow$ consensus
- Arachne
- Public assembler (mouse, several fungi)
- Overlap $\rightarrow$ layout $\rightarrow$ consensus
- Phusion
- Overlap $\rightarrow$ clustering $\rightarrow$ PHRAP $\rightarrow$ assemblage $\rightarrow$ consensus
- Euler
- Indexing $\rightarrow$ Euler graph $\rightarrow$ layout by picking paths $\rightarrow$ consensus


## EULER Fragment Assembly



- Traditional "overlap-layout-consensus" technique has a high rate of mis-assembly
- EULER uses the Eulerian Path approach borrowed from the SBH problem
- Fragment assembly without repeat masking can be done in linear time with greater accuracy


## Overlap Graph: Hamiltonian Approach



Each vertex represents a read from the original sequence.
Vertices from repeats are connected to many others.


Find a path visiting every VERTEX exactly once: Hamiltonian path problem

## Overlap Graph: Eulerian Approach





Placing each repeat edge together gives a clear progression of the path through the entire sequence.

Find a path visiting every EDGE exactly once:
Eulerian path problem

## Multiple Repeats


Repeat1 Repeat2 Repeat1 Repeat2



Can be easily constructed with any number of repeats


## Construction of Repeat Graph



- Construction of repeat graph from $k-$ mers: emulates an SBH experiment with a huge (virtual) DNA chip.
- Breaking reads into $k-$ mers: Transform sequencing data into virtual DNA chip data.


## Construction of Repeat Graph (cont'd)



- Error correction in reads: "consensus first" approach to fragment assembly. Makes reads (almost) error-free BEFORE the assembly even starts.
- Using reads and mate-pairs to simplify the repeat graph (Eulerian Superpath Problem).


## Approaches to Fragment Assembly



## Find a path visiting every VERTEX exactly once in the OVERLAP graph:

## Hamiltonian path problem



NP-complete: algorithms unknown

## Approaches to Fragment Assembly (cont'd) <br> 

Find a path visiting every EDGE exactly once in the REPEAT graph:

## Eulerian path problem



Linear time algorithms are known

## Making Repeat Graph Without DNA



- Problem: Construct the repeat graph from a collection of reads.

- Solution: Break the reads into smaller pieces.
- Virtual DNA chip allows the biological problem to be solved within the technological constraints.



## Repeat Sequences: Emulating a <br> DNA Chip (cont'd)

- Reads are constructed from an original sequence in lengths that allow biologists a high level of certainty.
- They are then broken again to allow the technology to sequence each within a reasonable array.


## Minimizing Errors



- If an error exists in one of the 20-mer reads, the error will be perpetuated among all of the smaller pieces broken from that read.



## Minimizing Errors (cont'd)



- However, that error will not be present in the other instances of the 20-mer read.
- So it is possible to eliminate most point mutation errors before reconstructing the original sequence.


## Conclusions



- Graph theory is a vital tool for solving biological problems
- Wide range of applications, including sequencing, motif finding, protein networks, and many more


## References



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http://www.mimg.ucla.edu/bobs/C159/Presentations/Benzer.pdf
- Batzoglou, S. Computational Genomics Course, Stanford University (2006). http://ai.stanford.edu/~serafim/CS262_2006/

